

REAL SCHUBERT CALCULUS: POLYNOMIAL SYSTEMS AND A CONJECTURE OF SHAPIRO AND SHAPIRO

FRANK SOTTILE

ABSTRACT. Boris Shapiro and Michael Shapiro have a conjecture concerning the Schubert calculus and real enumerative geometry and which would give infinitely many families of zero-dimensional systems of real polynomials (including families of overdetermined systems)—all of whose solutions are real. It has connections to the pole placement problem in linear systems theory and to totally positive matrices. We give compelling computational evidence for its validity, prove it for infinitely many families of enumerative problems, show how a simple version implies more general versions, and present a counterexample to a general version of their conjecture.

This is a companion paper to [15] and [47], which describe the mathematics involved in two spectacular computations verifying specific instances of this conjecture.

1. INTRODUCTION

Determining the number of real solutions to a system of polynomial equations is a challenging problem in symbolic and numeric computation [18, 44] with real world applications. Related questions include when a problem of enumerative geometry can have all solutions real [39] and when may a given physical system be controlled by real output feedback [6, 32, 45]. In May 1995, Boris Shapiro and Michael Shapiro communicated to the author a remarkable conjecture connecting these three lines of inquiry.

They conjectured a relation between topological invariants of the real and of the complex points in an intersection of Schubert cells in a flag manifold, if the cells are chosen according to a recipe they give. When the intersection is zero-dimensional, this asserts that all points are real. Their conjecture is false—we give full description and present a counterexample in Section 5. However, there is considerable evidence for their conjecture if the Schubert cells are in a Grassmann manifold. It is this variant which is related to the lines of inquiry above and which this paper is about.

Here is the simplest (but still very interesting and open) special case of this conjecture: Let $m, p > 1$ be integers and let X be a $p \times m$ -matrix of indeterminates. Let

1991 *Mathematics Subject Classification.* 12D10, 13P10, 14N10, 14M15, 14Q20, 14P99, 15A48, 93B55.

Key words and phrases. Enumerative geometry, Grassmannian, Gröbner basis, Overdetermined system.

Supported in part by NSERC grant OGP0170279 and NSF grants DMS-9022140 and DMS-9701755.

$K(s)$ be the $m \times (m + p)$ -matrix of polynomials in s whose i, j th entry is

$$(1) \quad \binom{j-i}{i-1} s^{j-i}.$$

Set

$$\varphi_{m,p}(s; X) := \det \begin{bmatrix} K(s) \\ I_p & X \end{bmatrix},$$

where I_p is the $p \times p$ identity matrix.

Conjecture 1.1 (Shapiro-Shapiro). *For all integers $m, p > 1$, the polynomial system*

$$(2) \quad \varphi_{m,p}(1; X) = \varphi_{m,p}(2; X) = \cdots = \varphi_{m,p}(mp; X) = 0$$

is zero-dimensional with

$$(3) \quad d_{m,p} := \frac{1! 2! 3! \cdots (p-2)! (p-1)! \cdot (mp)!}{m! (m+1)! (m+2)! \cdots (m+p-1)!}$$

solutions, and all of them are real.

It is a Theorem of Schubert [36] that $d_{m,p}$ is a sharp bound for the number of isolated solutions. Conjecture 1.1 has been verified for all $1 < m \leq p$ with $mp \leq 12$. The case of $(m, p) = (3, 4)$ ($d_{m,p} = 462$) is due to an heroic calculation of Faugère, Rouillier, and Zimmermann [15] (see Section 2.4 for a discussion).

Conjecture 1.1 is related to a question of Fulton [16, §7.2]: “How many solutions to a problem in enumerative geometry may be real, where that problem consists of counting figures of some kind having a given position with respect to some given (fixed) figures.”. For 2-planes having a given position with respect to fixed linear subspaces, the answer is that all may be real [40]. This was also shown for the problem of 3264 plane conics tangent to five given conics [31]. More examples, including that of 3-planes in \mathbb{C}^6 meeting 9 given 3-planes nontrivially, are found in [39, 41]. The result [15] extends this to 3-planes in \mathbb{C}^7 meeting 12 given 4-planes nontrivially.

Only the simplest form of the conjecture of Shapiro and Shapiro has appeared in print [23, 33, 39]. While more general forms have circulated informally, there is no definitive source describing the conjectures or the compelling evidence that has accumulated (or a counterexample to the original conjecture). The primary aim of this paper is to rectify this situation and make these conjectures available to a wider audience.

In Section 2, we describe a version of the conjecture related to the pole placement problem of linear systems theory. For this, the integers $1, 2, \dots, mp$ in the polynomial system (2) of Conjecture 1.1 are replaced by generic real numbers and all $d_{m,p}$ solutions are asserted to be real. We present evidence (computational and Theorems) in support of it. Subsequent sections describe the conjecture in greater generality—for enumerative problems arising from the Schubert calculus on Grassmannians in Section 3 and a newer extension involving totally positive matrices [1] in Section 4. We describe and give evidence for each extension and show how the version of the

conjecture in Section 2 implies more general versions involving Pieri-type enumerative problems. In Section 5, we present a counterexample to their original conjecture and discuss further questions.

A remark on the form of these conjectures is warranted. Conjecture 1.1 gives an infinite list of specific polynomial systems, and conjectures that each has only real solutions. The full conjectures are richer. For each collection of *Schubert data*, Shapiro and Shapiro give a continuous family of polynomial systems and conjecture that each of the resulting systems of polynomials has only real solutions. Conjecture 1.1 concerns one specific polynomial system in each family, for an infinite subset of Schubert data.

Results here were aided or are due to computations. Further documentation including Maple V.5 and Singular 1.2.1 [19] scripts used are available on the web page [42].

CONTENTS

1. Introduction	1
2. Linear equations in Plücker coordinates	3
2.1. Some enumerative geometry	3
2.2. The conjecture of Shapiro and Shapiro	4
2.3. Pole placement problem	5
2.4. Computational evidence	6
2.5. Equivalent systems	8
2.6. Conjecture 2.1 for $m = 2, p = 3$	10
3. Schubert conditions on a Grassmannian	12
3.1. The Schubert calculus on $Grass(p, m + p)$	12
3.2. Systems of polynomials	13
3.3. Pieri Schubert conditions	14
3.4. An infinite family	17
3.5. Computational evidence	19
4. Total positivity	22
5. Further remarks	25
5.1. A counterexample to the original conjecture	25
5.2. Further questions	26
Acknowledgements	28
References	28

2. LINEAR EQUATIONS IN PLÜCKER COORDINATES

2.1. Some enumerative geometry. Consider the following problem in enumerative geometry: How many p -planes meet mp general m -planes in \mathbb{C}^{m+p} nontrivially?

The set of p -planes in \mathbb{C}^{m+p} , $Grass(p, m + p)$, is called the *Grassmannian of p -planes in \mathbb{C}^{m+p}* . This complex manifold of dimension mp is an algebraic subvariety of the projective space $\mathbb{P}^{\binom{m+p}{p}-1}$. To see this, represent a p -plane X in \mathbb{C}^{m+p} as the row space of a $p \times (m + p)$ -matrix, also written X . The maximal minors of X are

its *Plücker coordinates* and determine a point in $\mathbb{P}^{\binom{m+p}{p}-1}$. This gives the Plücker embedding of $Grass(p, m+p)$. If X is generic, then its first p columns are linearly independent, so we may assume they form a $p \times p$ -identity matrix. The remaining mp entries determine X uniquely and give local coordinates for $Grass(p, m+p)$, showing it has dimension mp .

Consider a m -plane K to be the row space of a $m \times (m+p)$ -matrix, also written K . Then $K \cap X$ is nontrivial if and only if

$$\det \begin{bmatrix} K \\ X \end{bmatrix} = 0.$$

Laplace expansion along X gives a linear equation in the Plücker coordinates of X .

If K_1, \dots, K_{mp} are m -planes in general position, then the conditions that X meet each of the K_i nontrivially are mp linear equations in the Plücker coordinates of X , and these are independent by Kleiman's Transversality Theorem [24]. Hence there are finitely many p -planes X which meet each K_i nontrivially and this number is the degree of $Grass(p, m+p)$ in $\mathbb{P}^{\binom{m+p}{p}-1}$, which Schubert [36] determined to be $d_{m,p}$.

2.2. The conjecture of Shapiro and Shapiro. Shapiro and Shapiro gave a recipe for selecting real m -planes K_1, \dots, K_{mp} and conjecture that when they are in general position, all $d_{m,p}$ p -planes meeting each K_i are real. The standard rational normal curve is the image of the map $\gamma : \mathbb{R} \rightarrow \mathbb{R}^{m+p}$ given by

$$(4) \quad \gamma : s \longmapsto (1, s, s^2, \dots, s^{m+p-1})$$

Then the matrix $K(s)$ of the Introduction (1) has rows

$$\gamma(s), \gamma'(s), \frac{\gamma''(s)}{2}, \dots, \frac{\gamma^{(m-1)}(s)}{(m-1)!},$$

where we take derivatives with respect to the parameter s . Let X be a $p \times m$ -matrix of indeterminates. Define

$$\varphi_{m,p}(s; X) := \det \begin{bmatrix} K(s) \\ I_p \ X \end{bmatrix}.$$

Conjecture 2.1 (Shapiro-Shapiro). *For all integers $m, p > 1$ and almost all distinct real numbers s_1, \dots, s_{mp} , the system of mp equations*

$$(5) \quad \varphi_{m,p}(s_1; X) = \varphi_{m,p}(s_2; X) = \dots = \varphi_{m,p}(s_{mp}; X) = 0$$

is zero-dimensional with $d_{m,p}$ real solutions.

Let $K(s)$ denote both the $m \times (m+p)$ -matrix defined above and its row space, an m -plane. Conjecture 2.1 asserts that the m -planes $K(s_1), \dots, K(s_{mp})$ are in general position, and any p -plane meeting each $K(s_i)$ is real. The systems are zero-dimensional [3, 11] and there are generically no multiplicities. Conjecture 1.1 is the special case when $s_i = i$.

Example 2.2. We establish Conjecture 2.1 when $m = p = 2$. Then

$$\varphi_{2,2}(s; X) = \det \begin{bmatrix} 1 & s & s^2 & s^3 \\ 0 & 1 & 2s & 3s^2 \\ 1 & 0 & x_{11} & x_{12} \\ 0 & 1 & x_{21} & x_{22} \end{bmatrix}$$

is

$$s^4 - 2s^3x_{21} + s^2x_{22} - 3s^2x_{11} + 2sx_{12} + x_{11}x_{22} - x_{12}x_{21}.$$

We show that if $s, t, u, v \in \mathbb{R}$ are distinct, then the system of polynomial equations

$$(6) \quad \varphi_{2,2}(s) = \varphi_{2,2}(t) = \varphi_{2,2}(u) = \varphi_{2,2}(v) = 0$$

has all $d_{2,2} = 2$ solutions real. Our method will be to solve (6) by elimination.

Let e_i be the i th elementary symmetric polynomial in s, t, u, v . In the lexicographic term order with $x_{11} > x_{12} > x_{22} > x_{21}$ on the ring $\mathbb{Q}(s, t, u, v)[x_{11}, x_{12}, x_{22}, x_{21}]$, the ideal $\langle \varphi_{2,2}(s), \varphi_{2,2}(t), \varphi_{2,2}(u), \varphi_{2,2}(v) \rangle$ has a Gröbner basis consisting of the following polynomials:

$$2x_{21} - e_1, \quad x_{22} - 3x_{11} - e_2, \quad 2x_{12} + e_3, \quad \text{and} \quad 12x_{11}^2 + 4e_2x_{11} + e_1e_3 - 4e_4.$$

Thus, for distinct s, t, u, v , the system (6) has 2 solutions and they are real if the discriminant of the last equation,

$$16e_2^2 - 48e_1e_3 + 192e_4,$$

is positive. Expanding this discriminant in the parameters s, t, u, v , we obtain

$$8((s-t)^2(u-v)^2 + (s-u)^2(t-v)^2 + (s-v)^2(t-u)^2).$$

Hence all solutions are real, establishing Conjecture 2.1 when $m = p = 2$. Theorem 2.3 proves Conjecture 2.1 when $(m, p) = (2, 3)$.

2.3. Pole placement problem. Suppose we have a physical system (for example, a mechanical linkage) with inputs $u \in \mathbb{R}^m$ and outputs $y \in \mathbb{R}^p$ for which there are internal states $x \in \mathbb{R}^n$ such that the system evolves by the first order linear differential equation

$$(7) \quad \begin{aligned} \dot{x} &= Ax + Bu, \\ y &= Cx. \end{aligned}$$

(We assume n is the minimal number of internal states needed to obtain a first order equation.) If the input is controlled by constant output feedback, $u = Xy$, then we obtain

$$\dot{x} = (A + BXC)x.$$

The natural frequencies of this controlled system are the roots s_1, \dots, s_n of

$$(8) \quad \varphi(s) := \det(sI_n - A - BXC).$$

The pole assignment problem asks the inverse question: Given a system (7) and a polynomial $\varphi(s)$ of degree n , which feedback laws X satisfy (8)?

A coprime factorization of the transfer function is two matrices $N(s)$, $D(s)$ of polynomials with $\det(D(s)) = \det(sI_n - A)$ and $N(s)D(s)^{-1} = C(sI_n - A)^{-1}B$. This always exists. A standard transformation (*cf.* [6, §2]) shows that, up to a sign of ± 1 ,

$$(9) \quad \varphi(s) = \det \begin{bmatrix} N(s) & D(s) \\ I_p & X \end{bmatrix},$$

If we set $K(s) := [N(s) \ D(s)]$, write $K(s)$ for the m -dimensional row space of this matrix, and let X be the p -plane $[I_p \ X]$, then (9) is equivalent to

$$(10) \quad X \cap K(s_i) \neq \{0\} \quad \text{for } i = 1, \dots, n,$$

where s_1, \dots, s_n are the roots of $\varphi(s)$.

If the m -planes $K(s_1), \dots, K(s_n)$ are in general position, then $mp \geq n$ is necessary for there to be any feedback laws. These m -planes are not *a priori* in general position.

To see this, let $K : \mathbb{P}^1 \rightarrow \text{Grass}(m, m+p)$ be the extension of the map given by $s \mapsto K(s)$. Then K is a parameterized rational curve of degree n in $\text{Grass}(m, m+p)$. The space of all such curves K with n distinguished points $\{K(s_1), \dots, K(s_n)\}$ has dimension [43]

$$mp + n(m+p) + n.$$

The space of all n -tuples of m -planes has dimension nmp . Therefore when

$$n > mp/(mp - m - p - 1),$$

such n -tuples constitute a proper subvariety of all n -tuples of m -planes.

However, the General Position Lemma [6] (see also [11]) states that there is a Zariski open subset of the data A, B, C, φ such that the m -planes $K(s_1), \dots, K(s_n)$ are in general position in that the set of X satisfying (10) has dimension $mp - n$.

Since all rational curves $K : \mathbb{P}^1 \rightarrow \text{Grass}(p, m+p)$ of degree n with $K(\infty) = [0 \ I_p]$ arise in this way [28], the polynomial systems of Conjecture 2.1 are instances of the pole placement problem. Interestingly, these very systems figure prominently in a proof of the General Position Lemma [4].

An important question is whether a given real system may be controlled by *real* feedback [5, 32, 33, 45, 48]: If all roots of $\varphi(s)$ are real, are there any real feedback laws X satisfying (9)? Few specific examples have been computed [7, 29, 33, 48]. In [33] an attempt was made to gauge how likely it is for a real system to be controllable by real feedback and how many of the feedback laws are real—in the case of $(m, p) = (2, 4)$ so that $d_{m,p} = 14$. In all, 600 different curves $K(s)$ were generated, and each of these were combined with 25 polynomials $\varphi(s)$ having 8 real roots. Only 7 of the resulting 15,000 systems had all feedback laws real. This is in striking contrast to the systems given in Conjecture 2.1, where all the feedback laws are conjectured to be real.

2.4. Computational evidence. Consider (9) as a map $\text{Grass}(p, m+p) \rightarrow \mathbb{P}^{mp}$ in local coordinates which associates a p -plane X to a polynomial φ (modulo scalars) of degree at most mp . When $K(s)$ is the curve $K_{m,p}(s)$ of Conjecture 2.1, the inverse image of the polynomial 1 is the single real point $[0 \ I_p]$. Rosenthal suggested that the fibre over a nearby polynomial may consist of $d_{m,p}$ real points.

Inspired by this, Rosenthal and Sottile [33] tested and verified several thousand instances of Conjecture 2.1 when $(m, p) = (2, 4)$. Each was a specific choice of m, p and mp distinct real numbers s_1, \dots, s_{mp} for which we showed all solutions to (5) are real. Any verified instance implies that all nearby instances in the space of parameters s_1, \dots, s_{mp} has all of its solutions real. In light of the computations described in Section 2.3, we felt this provided overwhelming evidence for the validity of Conjecture 2.1.

Our method was to solve the polynomial systems by elimination (see [10, §2] for a discussion of methods to solve systems of polynomial equations). We first choose distinct integral values of the parameters s_i and generate the resulting system of integral polynomial equations. Since we are performing an exact symbolic computation, we necessarily work with integral polynomials. Next, we compute an eliminant, a univariate polynomial $g(x)$ with the property that its roots are the set of x -coordinates of solutions to our system. When $g(x)$ has $d = d_{m,p}$ roots (Schubert's bound), there is a lexicographic Gröbner basis satisfying the Shape Lemma, since this system is zero-dimensional [11]. It follows that the solutions are rational functions (quotients of integral polynomials) of the roots of $g(x)$. In some instances, the eliminant we calculated did not have d roots. For these we found a different eliminant with d roots. Lastly, we checked that these eliminants had only real roots.

Table 1 gives the number of instances we know have been checked. By Lemma 3.7(ii), there is a bijection between instances of $(m, p) = (a, b)$ and $(m, p) = (b, a)$. Table 1 also lists the running time to compute a degree reverse lexicographic Gröbner basis for the systems of Conjecture 1.1, and the size of that basis. This used Singular-1.2.1 [19] on a K6-2-300 processor with 256M running Linux. The instances reported in the last 3 columns are not due the the author. A more complete account is found in [42].

m, p	4,2	5,2	3,3	6,2	7,2	4,3	2,8
$d_{m,p}$	14	42	42	132	429	462	1430
# checked	> 12,000	1000	550	55	2	2	1
time (sec)	.04	1.42	1.50	78.6	8175	—	—
size	1.4K	12.8K	18.6K	202K	4.58M	32M	—

TABLE 1. Instances checked

The computations of the last 2 columns stand out. The first is the case 8, 2 (also one instance each of 7, 2 and 4, 3) computed by Jan Verschelde [47] using his implementation of the SAGBI homotopy algorithm in [23]. Since the polynomial system of Conjecture 2.1 was ill-conditioned, he instead used the equivalent system of Conjecture 2.1' (in Section 2.5 below), where the $P_i(s)$ were the Chebyshev polynomials. These numerical calculations give approximate solutions whose condition numbers determine a neighborhood containing a solution. The solutions of this real system are stable under complex conjugation, so it sufficed to check that each neighbourhood and its complex conjugate were disjoint from all other neighborhoods. This computation took approximately 25 hours on a 166MHz Pentium II processor with 64M

running Linux. These algorithms are ‘embarrassingly parallelizable’, and in principle they can be used to check far larger polynomial systems.

The second is the case of $(m, p) = (3, 4)$ of Conjecture 1.1 (also all smaller cases with $m \leq p$), computed by Faugère, Rouillier, and Zimmermann [15]. They first used FGB [13] to calculate a degree reverse lexicographic Gröbner basis for the system (2) for $(m, p) = (3, 4)$ with $s_i = i$. This yielded a Gröbner basis of size 32M. They then computed a rational univariate representation [35] (a sophisticated substitute for an eliminant) in two ways. Once using a multi-modular implementation of the FGLM [14] algorithm and a second time using RS, an improvement of the RealSolving software [34] under development. The eliminant had degree 462 and size 3M, thus its general coefficient had 2,000 digits. Using an early implementation of Uspensky’s algorithm, they verified that all of its zeroes were real, proving Conjecture 1.1 for $(m, p) = (3, 4)$. In the course of this calculation, they found it necessary to rewrite their software.

2.5. Equivalent systems. The extension of the map (4) to \mathbb{P}^1

$$\gamma : [t, s] \longmapsto [t^{m+p-1}, st^{m+p-2}, \dots, s^{m+p-2}t, s^{m+p-1}]$$

is a parameterization of the standard real rational normal curve in \mathbb{P}^{m+p-1} and $K(s)$ is the m -plane osculating this curve at the point $\gamma[1, s]$. In general, a parameterized real rational normal curve is a map $\gamma : \mathbb{P}^1 \rightarrow \mathbb{P}^{m+p-1}$ of the form

$$[t, s] \longmapsto [P_1(s, t), P_2(s, t), \dots, P_{m+p}(s, t)]$$

where $P_1(s, t), \dots, P_{m+p}(s, t)$ form a basis for the space of real homogeneous polynomials in s, t of degree $m + p - 1$. All parameterized real rational normal curves are conjugate by a real projective transformation of \mathbb{P}^{m+p-1} . We deduce

Conjecture 2.1 (Geometric form) *For all integers $m, p > 1$ and almost all choices of mp m -planes K_1, \dots, K_{mp} osculating a real rational normal curve at distinct real points, there are exactly $d_{m,p}$ p -planes X satisfying*

$$X \cap K_i \neq \{0\} \quad \text{for } i = 1, \dots, mp$$

and all of these p -planes X are real.

Thus Conjecture 2.1 is equivalent to a conjecture concerning a much richer class of polynomial systems.

Conjecture 2.1’. *Suppose $m, p > 1$ are integers and $P_1(s), \dots, P_{m+p}(s)$ are a basis of the space of real polynomials of degree at most $m + p - 1$. Let $K(s)$ be the $m \times (m + p)$ matrix of polynomials whose i, j th entry is $P_j^{(i-1)}(s)$. Set*

$$\varphi(s; X) := \det \begin{bmatrix} K(s) \\ I_p X \end{bmatrix}.$$

Then, for almost all choices of distinct real numbers s_1, \dots, s_{mp} , the system

$$\varphi(s_1; X) = \varphi(s_2; X) = \dots = \varphi(s_{mp}; X) = 0$$

has exactly $d_{m,p}$ solutions, and all of them are real.

The polynomial matrix $K(s)$ of Conjecture 2.1' differs from that of Conjecture 2.1 by right multiplication by an invertible $(m+p) \times (m+p)$ -matrix. Thus the resulting polynomial systems differ primarily by choice of local coordinates for the Grassmannian. In linear systems theory, two physical systems are output feedback-equivalent if their matrices of coprime factors $[N(s) \ D(s)]$ differ in this manner [30].

We give an equivalent conjecture concerning a simpler system of polynomials with 2 fewer equations and unknowns. We may reparameterize the curve $K(s)$ of Conjecture 2.1 and assume $s_{mp-1} = 0$ and $s_{mp} = \infty$. Observe that $K(0) = [I_p \ 0]$ and $K(\infty) = [0 \ I_p]$. The collection of all p -planes X satisfying

$$(11) \quad X \cap [I_p \ 0] \neq \{0\} \quad \text{and} \quad X \cap [0 \ I_p] \neq \{0\}$$

is an irreducible rational variety of dimension $mp - 2$.

Let \mathcal{X} be the set of all $p \times (m+p)$ -matrices X whose entries $x_{i,j}$ satisfy:

$$x_{i,j} = 1 \quad \text{if} \quad j = i < p \quad \text{or} \quad (i,j) = (p, p+1)$$

$$x_{i,j} = 0 \quad \text{if} \quad \begin{cases} i = 1 & \text{and} \quad j \geq m \\ 1 < i < p & \text{and} \quad j < i \text{ or } j > i + m \\ i = p & \text{and} \quad j \leq p \end{cases}$$

The remaining $mp - 2$ entries are unconstrained and give coordinates for \mathcal{X} . The row space of a matrix X is a p -plane X satisfying (11) and almost all such p -planes arise in this fashion. Thus \mathcal{X} parameterizes a dense subset of the subvariety of p -planes X satisfying (11).

For example, if $(m, p) = (4, 3)$, then \mathcal{X} is the set of all matrices of the form:

$$\begin{bmatrix} 1 & x_{12} & x_{13} & x_{14} & 0 & 0 & 0 \\ 0 & 1 & x_{23} & x_{24} & x_{25} & x_{26} & 0 \\ 0 & 0 & 0 & 1 & x_{35} & x_{36} & x_{37} \end{bmatrix}.$$

Since the 1, m th entry of a matrix in \mathcal{X} vanishes,

$$\det \begin{bmatrix} K(s) \\ X \end{bmatrix}$$

factors as $s \cdot \psi(s; X)$.

Conjecture 2.1''. *Let $m, p > 1$ be integers. Then, for almost all choices of non-zero real numbers s_1, \dots, s_{mp-2} , the system of equations*

$$(12) \quad \psi(s_1; X) = \psi(s_2; X) = \dots = \psi(s_{mp-2}; X) = 0$$

is zero-dimensional with $d_{m,p}$ solutions, and all of them are real.

The systems of Conjecture 2.1 and the variations given here are deficient: They have fewer solutions than standard combinatorial bounds. For example, the system (12) consists of $mp - 2$ equations of degree p , thus its Bézout number is p^{mp-2} . A better combinatorial bound is the normalized volume of the Newton polytope $\mathcal{A}_{m,p}$ of the polynomial ψ [26]. Table 2 compares these combinatorial bounds with $d_{m,p}$, for some values of m, p . The volumes of $\mathcal{A}_{m,p}$ were computed using PHC [46], a software package for performing general polyhedral homotopy continuation. Note the striking difference between the equivalent systems m, p and p, m .

$m, p:$	2,2	3,2	4,2	5,2	6,2	7,2	8,2	2,3	3,3	4,3	2,4	3,4
$d_{m,p}:$	2	5	14	42	132	429	1430	5	42	462	14	462
$\text{vol } \mathcal{A}_{m,p}:$	2	5	18	67	248	919	3426	5	130	3004	42	7156
$p^{mp-2}:$	4	16	64	256	1024	4096	16384	81	2187	59,049	4096	1048576

TABLE 2. Combinatorial bounds vs. $d_{m,p}$

2.6. Conjecture 2.1 for $m = 2, p = 3$.

Theorem 2.3. *Conjecture 2.1 holds for $(m, p) = (2, 3)$.*

L. Gonzalez-Vega has also obtained this using resultants and Sturm-Habicht sequences.

Proof. We will prove the equivalent Conjecture 2.1''. Let $X := \{x_{12}, x_{23}, x_{24}, x_{35}\}$ be indeterminates. Set

$$\psi(s; X) = \det \begin{bmatrix} 1 & s & s^2 & s^3 & s^4 \\ 0 & 1 & 2s & 3s^2 & 4s^3 \\ 1 & x_{12} & 0 & 0 & 0 \\ 0 & 1 & x_{23} & x_{24} & 0 \\ 0 & 0 & 0 & 1 & x_{35} \end{bmatrix}$$

We solve the system of polynomials

$$(13) \quad \psi(s; X) = \psi(t; X) = \psi(u; X) = \psi(v; X) = 0$$

by elimination.

The ideal $\langle \psi(s), \psi(t), \psi(u), \psi(v) \rangle$ in the ring $\mathbb{Q}(s, t, u, v)[x_{12}, x_{23}, x_{24}, x_{35}]$ has degree $5 = d_{2,3}$ and the lexicographic Gröbner basis with $x_{12} < x_{23} < x_{24} < x_{35}$ contains the following univariate polynomial g , which is the universal eliminant for this family of systems:

$$x_{35}^5 - 4e_1x_{35}^4 + (4e_1^2 + 6e_2)x_{35}^3 - (12e_1e_2 + 4e_3)x_{35}^2 + (9e_2^2 + 8e_1e_3 - 4e_4)x_{35} - (12e_2e_3 - 8e_1e_4)$$

Here e_i is the i th the elementary symmetric polynomial in s, t, u, v . We show that g has 5 distinct real roots for every choice of distinct parameters s, t, u, v . The discriminant Δ of g has degree 20 in the variables s, t, u, v and 711 terms:

$$\begin{aligned} & 9e_3^4e_2^4e_1^4 - 54e_3^4e_2^3e_1^2 + 81e_3^4e_2^4 - 32e_3^5e_1^5 + 204e_3^5e_2e_1^3 - 324e_3^5e_2^2e_1 - 108e_3^6e_1^2 + 324e_3^6e_2 \\ & + 81e_4^2e_2^4e_1^4 - 486e_4^2e_2^5e_1^2 + 729e_4^2e_2^6 - 54e_4e_3^2e_2^3e_1^4 + 324e_4e_3^2e_2^4e_1^2 - 486e_4e_3^2e_2^5 \\ & + 204e_4e_3^3e_2e_1^5 - 1296e_4e_3^3e_2^2e_1^3 + 2052e_4e_3^3e_2^3e_1 - 8e_4e_3^4e_1^4 + 738e_4e_3^4e_2e_1^2 - 2106e_4e_3^4e_2^2 \\ & - 108e_4e_3^5e_1 - 324e_4^2e_3e_2^2e_1^5 + 2052e_4^2e_3e_2^3e_1^3 - 3240e_4^2e_3e_2^4e_1 - 108e_4^2e_3^2e_1^6 + 738e_4^2e_3^2e_2e_1^4 \\ & - 2592e_4^2e_3^2e_2^2e_1^2 + 3834e_4^2e_3^2e_2^3 - 368e_4^2e_3^3e_1^3 + 1800e_4^2e_3^3e_2e_1 - 27e_4^2e_3^4 + 324e_4^3e_2e_1^6 \\ & - 2106e_4^3e_2^2e_1^4 + 3834e_4^3e_2^3e_1^2 - 972e_4^3e_2^4 - 108e_4^3e_3e_1^5 + 1800e_4^3e_3e_2e_1^3 - 5544e_4^3e_3e_2^2e_1 \\ & - 634e_4^3e_3^2e_1^2 + 984e_4^3e_3^2e_2 - 27e_4^4e_1^4 + 984e_4^4e_2e_1^2 + 432e_4^4e_2^2 - 352e_4^4e_3e_1 - 64e_4^5. \end{aligned}$$

This vanishes when g has a double root. Thus the number of real roots of g is constant on each connected component (in \mathbb{R}^4) of the locus $\Delta \neq 0$. We show there is only one

connected component, and so the number of real roots of g (and thus the original system) does not depend upon the choice of real parameters. Since the roots of g evaluated at $(s, t, u, v) = (1, 2, 3, 4)$ are

$$8, 8 \pm \sqrt{19}, 8 \pm \sqrt{11}$$

it follows that there are always five real roots of g , and thus the system (13) has $d_{2,3} = 5$ real solutions whenever s, t, u, v are real and distinct.

We complete the proof. For $w \in \mathbb{Z}_{\geq 0}^{10}$, consider the polynomial:

$$(14) \quad s^{w_1} t^{w_2} u^{w_3} v^{w_4} (s-t)^{w_5} (s-u)^{w_6} (s-v)^{w_7} (t-u)^{w_8} (t-v)^{w_9} (u-v)^{w_{10}}$$

Let A_w be the primitive part of the symmetrization of this polynomial. Thus A_w is a sum of squares, none of which vanish on the locus where s, t, u, v are distinct. Then Δ is

$$\begin{aligned} \frac{1}{2} & (7A_{2220222222} + 3A_{2222402204} + 6A_{4222022222} + 7A_{4220222222} + 2A_{4420022222} \\ & + 2A_{2222440022} + 2A_{0222443022} + A_{4420202222} + 2A_{4222420022} + A_{4220022422} \\ & + A_{0222442202} + A_{2202024422} + 6A_{2222420024} + 10A_{4220022242} + 3A_{2222222222}). \end{aligned}$$

Note that the term $7A_{4220222222}$ does not vanish when a single parameter is zero. Similarly, the term $3A_{2222402204}$ does not vanish when $s = u$ and $t = v$ (but $u \neq t$). Thus the locus where $\Delta = 0$ has dimension 1 and so its complement is connected. \blacksquare

We have a Maple program which performs the computations described and runs in ~ 15 seconds on a K6-2-300 processor.

A *positive semidefinite* polynomial is a real polynomial that takes only nonnegative values. In the proof we showed Δ is positive semidefinite by exhibiting it as a sum of squares. Not all positive semidefinite polynomials are sums of squares of polynomials. There exist positive semidefinite polynomials of degree l in k variables which are not sums of squares of polynomials if $\min(k, l) > 2$ and $(k, l) \neq (3, 4)$ [21]. For Δ , $(k, l) = (4, 20)$.

The form of the squares we used (14) for the discriminant Δ , while motivated by the observation that no two parameters $(0, s, t, u, v, \infty)$ should coincide, is justified by the observation that any real zero of Δ must also be a zero of all the squares, if Δ is a sum of squares. (See [8] for other applications of this idea.)

Each of the polynomials A_w is a sum of squares, the number given by the orbit of the symmetric group on its index w . Since 6 have trivial stabilizer, 7 are stabilized by a transposition, one by the dihedral group D_8 , and one is invariant, there are $6 \cdot 24 + 7 \cdot 12 + 3 + 1 = 232$ squares in all. This is not the best possible. Choi, Lam, and Reznick [9] show, for degree l homogeneous polynomials in k variables that are a sum of squares of polynomials, at most

$$\Lambda(k, l) := \left\lfloor \frac{1}{2} \left(\sqrt{1 + 8 \binom{k+l-1}{l}} - 1 \right) \right\rfloor$$

squares are needed. Note that $\Lambda(4, 20) = 59$.

3. SCHUBERT CONDITIONS ON A GRASSMANNIAN

3.1. The Schubert calculus on $Grass(p, m+p)$. The enumerative problems of Section 2 are special cases of more general problems given by Schubert conditions on $Grass(p, m+p)$. A *Schubert condition* on $Grass(p, m+p)$ is an increasing sequence of integers

$$\alpha : 1 \leq \alpha_1 < \alpha_2 < \cdots < \alpha_p \leq m+p.$$

Let $\binom{[m+p]}{p}$ be the set of all such sequences. A *Schubert variety* $\Omega_\alpha F_\bullet$ is given by a Schubert condition α and a complete flag K_\bullet in \mathbb{C}^{m+p} , a sequence of subspaces

$$K_\bullet : K_1 \subset K_2 \subset \cdots \subset K_{m+p} = \mathbb{C}^{m+p}$$

where $\dim K_i = i$. Then the Schubert variety $\Omega_\alpha F_\bullet$ is the set of all p -planes X satisfying

$$(15) \quad \dim X \cap K_{m+p+1-\alpha_i} \geq p+1-i$$

for each $i = 1, 2, \dots, p$. This irreducible subvariety of $Grass(p, m+p)$ has codimension $|\alpha| := \sum_i (\alpha_i - i)$.

A sequence $\alpha^\bullet = \alpha^1, \dots, \alpha^n$ with $\alpha^j \in \binom{[m+p]}{p}$ with $\sum_j |\alpha^j| = mp$ is *Schubert data* for $Grass(p, m+p)$. Given Schubert data α^\bullet and flags $K_\bullet^1, \dots, K_\bullet^n$ in general position, there are finitely many (complex) p -planes X which lie in the intersection of the Schubert varieties $\Omega_{\alpha^j} K_\bullet^j$ for $j = 1, \dots, n$. The classical Schubert calculus [25] gives the following recipe for computing this number $d = d(m, p; \alpha^\bullet)$. Let h_1, \dots, h_m be indeterminates with $\deg(h_i) = i$. For each integer sequence $\beta_1 < \beta_2 < \cdots < \beta_r$ define the following polynomial:

$$(16) \quad S_\beta := \det(h_{\beta_i - j})_{1 \leq i, j \leq r}.$$

Here $h_0 := 1$ and $h_i := 0$ if $i < 0$ or $i > m$. Let \mathcal{I} be the ideal in $\mathbf{Q}[h_1, \dots, h_m]$ generated by those S_β with $r = p+1$, $1 < \beta_1$, and $\beta_{p+1} \leq m+p$. The quotient ring $\mathcal{A}_{m,p} := \mathbf{Q}[h_1, \dots, h_m]/\mathcal{I}$ is isomorphic to the cohomology ring of $Grass(p, m+p)$. It is Artinian with one-dimensional socle in degree mp . In the socle we have the relation

$$d \cdot (h_m)^p - S_{\alpha^1} S_{\alpha^2} \cdots S_{\alpha^n} \in \mathcal{I}.$$

We can compute the number d by normal form reduction modulo any Gröbner basis for \mathcal{I} .

If γ is a rational normal curve, then the flag of subspaces osculating γ at a point is the *osculating flag* to γ at that point.

Conjecture 3.1. *Let $m, p > 1$ and α^\bullet be Schubert data for $Grass(p, m+p)$. For almost all choices of flags $K_\bullet^1, \dots, K_\bullet^n$ osculating a fixed rational normal curve at real points, there are exactly $d(m, p, \alpha^\bullet)$ p -planes X in the intersection of Schubert varieties*

$$\Omega_{\alpha^1} K_\bullet^1 \cap \Omega_{\alpha^2} K_\bullet^2 \cap \cdots \cap \Omega_{\alpha^n} K_\bullet^n,$$

and each of these p -planes is real.

As with Conjecture 2.1, the intersection is zero-dimensional if the points of osculation are distinct [11], and there are no multiplicities for the important class of Pieri Schubert data, (described below) which includes the case of Conjecture 2.1.

If $\alpha_i = 1 + \alpha_{i-1}$, then condition (15) for $i - 1$ implies (15) for i . Thus only those conditions (15) with $\alpha_i - \alpha_{i-1} > 1$ (or $\alpha_1 > 1$) are *essential*, and so only the subspaces $K_{m+p+1-\alpha_i}$ corresponding to essential conditions need be specified in a flag. If $\alpha := (1, 2, \dots, p-1, p+1)$, then only the last condition is essential, thus the Schubert variety $\Omega_\alpha K_\bullet$ consists of those X with $\dim X \cap K_m \geq 1$. This shows Conjecture 2.1 is a special case of Conjecture 3.1.

3.2. Systems of polynomials. A complete flag K_\bullet is represented by a nonsingular matrix also written K_\bullet : The i -plane K_i is the row space of K_i , the first i rows of K_\bullet . The condition that $\dim X \cap K_{m+p+1-\alpha_i} \geq p+1-i$ is given by

$$(m+p+1+i-\alpha_i)\text{-minors of } \begin{bmatrix} K_{m+p+1-\alpha_i} \\ X \end{bmatrix} = 0.$$

The flag $K_\bullet(s)$ osculating the rational normal curve γ with the parameterization (4) at $\gamma(s)$ is represented by the $(m+p) \times (m+p)$ -matrix whose i, j th entry is $\binom{j-i}{i-1} s^{j-i}$.

Conjecture 3.1'. *Let $m, p > 1$ and α^\bullet be Schubert data for $\text{Grass}(p, m+p)$. For almost all n -tuples of distinct real numbers s_1, \dots, s_n , the system of polynomials*

$$(m+p+1+i-\alpha_i^j)\text{-minors of } \begin{bmatrix} K_{m+p+1-\alpha_i^j}(s_j) \\ I_p \ X \end{bmatrix} = 0$$

for $i = 1, \dots, p$ and $j = 1, \dots, n$ has $d(m, p, \alpha^\bullet)$ solutions, and each is real.

For any Schubert conditions α, β with $\alpha_i + \beta_{p+1-i} \leq m+p$ for $i = 1, \dots, p$, let $\mathcal{X}_{\alpha, \beta}$ be the collection of all $p \times (m+p)$ -matrices X whose entries x_{ij} satisfy

$$(17) \quad \begin{array}{ll} x_{i, \alpha_i} = 1 & \text{for } i = 1, \dots, p \\ x_{i, j} = 0 & \text{if } j < \alpha_i \text{ or } j > m+p+1-\beta_{p+1-i} \end{array}$$

If $X \in \mathcal{X}_{\alpha, \beta}$, then the row space of X is a p -plane in the intersection $\Omega_\alpha K_\bullet(\infty) \cap \Omega_\beta K_\bullet(0)$. In this way, $\mathcal{X}_{\alpha, \beta}$ parameterizes a Zariski open subset of the set of all such p -planes. This parameterization can be used to obtain a system of equations simpler than, but equivalent to, the system of Conjecture 3.1'.

The map $\mathcal{X}_{\alpha, \beta} \rightarrow \text{Grass}(p, m+p)$ is not injective. For example, $\mathcal{X}_{123, 134}$ consists of all 3×7 -matrices of the form:

$$\begin{bmatrix} 1 & x_{12} & x_{13} & x_{14} & 0 & 0 & 0 \\ 0 & 1 & x_{23} & x_{24} & x_{25} & 0 & 0 \\ 0 & 0 & 1 & x_{34} & x_{35} & x_{36} & x_{37} \end{bmatrix}$$

Let r_1, r_2, r_3 be the rows of such a matrix. If $x_{36} = x_{37} = 0$, then for each $a \in \mathbb{C}$, the matrix with rows $r_1, r_2 + ar_3, r_3$ is in $\mathcal{X}_{123, 134}$, and these all have the same row space. Similarly, if $x_{25} = 0$, then the same is true of the matrices with rows $r_1 + ar_2, r_2, r_3$.

Let $\mathcal{X}_{\alpha,\beta}^\circ \subset \mathcal{X}_{\alpha,\beta}$ be the set of those matrices whose entries further satisfy

$$(18) \quad \begin{aligned} &\text{For each } i = 2, \dots, p, \text{ at least one } x_{ij} \neq 0, \text{ for } j \text{ satisfying} \\ &\beta_{p+1-i} \leq m + p + 1 - j < \beta_{p+2-i}. \end{aligned}$$

The map $\mathcal{X}_{\alpha,\beta}^\circ \rightarrow \text{Grass}(p, m+p)$ is injective. For $\mathcal{H}_{123,134}$ this condition is that $x_{25} \neq 0$ and $(x_{36}, x_{37}) \neq (0, 0)$.

3.3. Pieri Schubert conditions. If $\alpha \in \binom{[m+p]}{p}$ has $\alpha_{p-1} = p-1$ and $\alpha_p = p+a$, then the Schubert variety $\Omega_\alpha K_\bullet$ is

$$\{X \mid X \cap K_{m+1-a} \neq \{0\}\}.$$

We call such a Schubert condition a *Pieri condition* and denote it by J_a . *Pieri Schubert data* are Schubert data $\alpha^1, \dots, \alpha^n$ where at most 2 of the conditions α^i are *not* Pieri conditions. These include the Schubert data of Conjecture 2.1.

Proposition 3.2 (Theorem 9.1 in [11]). *If α^\bullet are Pieri Schubert data and the flags $K_\bullet^1, \dots, K_\bullet^n$ osculate a rational normal curve at general points, then the intersection of Schubert varieties*

$$\Omega_{\alpha^1} K_\bullet^1 \cap \Omega_{\alpha^2} K_\bullet^2 \cap \dots \cap \Omega_{\alpha^n} K_\bullet^n,$$

is transverse. In particular, there are no multiplicities.

Here is the main theorem of this section.

Theorem 3.3. *Let $a, b > 1$ and suppose that Conjecture 2.1 holds for this $(m, p) = (a, b)$. Then Conjecture 3.1 holds for any Pieri Schubert data for $\text{Grass}(p, m+p)$ with $(p, m) \leq (a, b)$ or $(p, m) \leq (b, a)$, in each coordinate.*

Remark 3.4. If the conclusion of Proposition 3.2 held for all Schubert data, then the proof we give of Theorem 3.3 would imply its conclusion for all Schubert data as well.

Pieri conditions are special because of Pieri's formula. For $\alpha, \beta \in \binom{[m+p]}{p}$ and $a > 0$, we write $\alpha <_a \beta$ if $|\alpha| + a = |\beta|$ and

$$\alpha_1 \leq \beta_1 < \alpha_2 \leq \beta_2 < \dots < \alpha_p \leq \beta_p.$$

Proposition 3.5 (Pieri's Formula). *Let $J_a := 1 < 2 < \dots < p-1 < p+a \in \binom{[m+p]}{p}$.*

(i) *In the cohomology ring $\mathcal{A}_{m,p}$ of $\text{Grass}(p, m+p)$, $S_{J_a} = h_a$ and*

$$S_\alpha \cdot S_{J_a} = \sum_{\alpha <_a \beta} S_\beta.$$

(ii) *If $K_\bullet(s)$ and $K_\bullet(t)$ are flags osculating a rational normal curve at points s and t , then*

$$\lim_{s \rightarrow t} (\Omega_\alpha K_\bullet(t) \cap \Omega_{J_a} K_\bullet(s)) = \sum_{\alpha <_a \beta} \Omega_\beta K_\bullet(t).$$

Here, the limit is taken as cycles. By this we mean that the sum is the fundamental cycle of the limit of the schemes $\Omega_\alpha K_\bullet(t) \cap \Omega_{J_a} K_\bullet(s)$ as s approaches t along the rational normal curve.

(iii) Suppose $\alpha^\bullet = \alpha^1, J_a, \alpha^2, \dots, \alpha^n$ are Schubert data. Then

$$d(m, p; \alpha^\bullet) = \sum_{\alpha <_a \beta} d(m, p; \beta, \alpha^2, \dots, \alpha^n).$$

Statement (i) is the usual statement of Pieri's formula [17, 22], Statement (ii) is Theorem 8.1 of [11], and Statement (iii) is a direct consequence of (i).

Proof. We deduce Theorem 3.3 from Lemma 3.7 below, which shows some simple dependencies between Conjecture 3.1 for different collections of Schubert data. Definition (15) implies that $\Omega_\beta K_\bullet \subset \Omega_\alpha K_\bullet$ if and only if $\alpha \leq \beta$ coordinatewise. In fact, $\Omega_\beta K_\bullet \cap \Omega_\alpha K_\bullet = \Omega_{\beta \vee \alpha} K_\bullet$, where $\beta \vee \alpha$ is the coordinatewise maximum of α and β . We make some definitions needed for the statement of Lemma 3.7.

Definition 3.6. Let $m, p > 1$ be integers.

1. For $\alpha \in \binom{[m+p]}{p}$ define $\alpha^\perp \in \binom{[m+p]}{m}$ to be the increasing sequence obtained from the numbers $\{1, 2, \dots, m+p\} \setminus \{\alpha_1, \dots, \alpha_p\}$. Given Schubert data α^\bullet for $\text{Grass}(p, m+p)$, set α^{\perp} to be $(\alpha^1)^\perp, \dots, (\alpha^n)^\perp$.
2. Suppose $p > 2$. For $\alpha \in \binom{[m+p-1]}{p-1}$ define $\alpha^+ \in \binom{[m+p]}{p}$ to be $1 < 1 + \alpha_1 < \dots < 1 + \alpha_{p-1}$. Given Schubert data α^\bullet for $\text{Grass}(p, m+p)$, set α^+ to be $(\alpha^1)^+, \dots, (\alpha^n)^+$.
3. Let \preceq be the partial order on Pieri Schubert data where we say that β^\bullet covers $\alpha^\bullet = \alpha^1, \dots, \alpha^n$ if one of the following holds

$$\begin{aligned} \beta^\bullet &= \beta, \alpha^3, \dots, \alpha^n \text{ with } \alpha^2 = J_a \text{ and } \alpha^1 <_a \beta \\ \beta^\bullet &= \alpha^1, \dots, \alpha^{n-2}, \beta \text{ with } \alpha^{n-1} = J_a \text{ and } \alpha^n <_a \beta. \end{aligned}$$

Lemma 3.7. Let $m, p > 1$ be integers.

- (i) If α^\bullet is Schubert data for $\text{Grass}(p, m+p)$, then α^{\perp} is Schubert data for $\text{Grass}(m, m+p)$. Moreover, Conjecture 3.1 holds for m, p, α^\bullet if and only if it holds for m, α^{\perp} .
- (ii) Suppose $p > 2$ and let $J_m := 1 < 2 < \dots < p-1 < p+m$. If α^\bullet is Schubert data for $\text{Grass}(p-1, m+p-1)$, then $\beta^\bullet := \alpha^+, J_m$ is Schubert data for $\text{Grass}(m, m+p)$. Moreover, Conjecture 3.1 holds for $m, p-1, \alpha^\bullet$ if and only if it holds for m, p, β^\bullet .
- (iii) Let $\alpha^\bullet, \beta^\bullet$ be Pieri Schubert data for $\text{Grass}(p, m+p)$ with $\alpha^\bullet \preceq \beta^\bullet$. If Conjecture 3.1 holds for α^\bullet for $\text{Grass}(p, m+p)$, then it holds for β^\bullet .

Proof of Theorem 3.3. First note that Conjecture 3.1 holds for Schubert data α^\bullet for $\text{Grass}(p, m+p)$ if and only if it holds for any rearrangement of the data α^\bullet . Suppose Conjecture 2.1 holds for $\text{Grass}(b, a+b)$. Let α^\bullet be Pieri Schubert data for $\text{Grass}(p, m+p)$ where $(m, p) \leq (a, b)$ or $(m, p) \leq (b, a)$ coordinatewise. Since $J_1^\perp = J_1$, Conjecture 2.1 holds also for $(m, p) = (b, a)$, by Lemma 3.7(i). Thus we may assume that $(m, p) \leq (a, b)$. By Lemma 3.7(ii), there exist Pieri Schubert data β^\bullet for $\text{Grass}(b, a+b)$ such that Conjecture 3.1 holds for α^\bullet if and only if it holds for β^\bullet . Finally, Theorem 3.3 follows from (iii) by noting that the Schubert data of Conjecture 2.1, namely $\alpha^1 = \dots = \alpha^{ab} = J_1$, is minimal among all Pieri Schubert data for $\text{Grass}(b, a+b)$. \blacksquare

Proof of Lemma 3.7. For (i), fix a real inner product on \mathbb{C}^{m+p} . Then the map $X \mapsto X^\perp$ gives an isomorphism between $\text{Grass}(p, m+p)$ and $\text{Grass}(m, m+p)$. Given a flag K_\bullet and an increasing sequence α , let K_\bullet^\perp be the flag of annihilators of the subspaces of K_\bullet . Then we have

$$X \in \Omega_\alpha K_\bullet \iff X^\perp \in \Omega_{\alpha^\perp} K_\bullet^\perp.$$

Furthermore, if $K_\bullet(s)$ is the flag of subspaces osculating a rational normal curve γ at a point $\gamma(s)$, then $(K_{m+p-1}(s))^\perp$ is a rational normal curve with $K_\bullet^\perp(s)$ its osculating flag. Thus Conjecture 3.1 for Schubert data α^\bullet for $\text{Grass}(p, m+p)$ is equivalent to Conjecture 3.1 for Schubert data α^\perp for $\text{Grass}(m, m+p)$.

For (ii), let γ be the rational curve (4) with $K_\bullet(s)$ as before. Then $X \in \Omega_{J_m} K_\bullet(\infty)$ if and only if $\langle \gamma(\infty) \rangle = K_1(\infty) \subset X$. Consider the projection $\pi : \mathbb{C}^{m+p} \twoheadrightarrow \mathbb{C}^{m+p-1}$ from the last coordinate $\gamma(\infty)$. If $X \in \Omega_{J_m} K_\bullet(\infty)$, then $X' := \pi X$ is a $(p-1)$ -plane. This induces an isomorphism $\pi : \Omega_{J_m} K_\bullet(\infty) \xrightarrow{\sim} \text{Grass}(p-1, m+p-1)$. The inverse map is given by $X' \mapsto K_1(\infty) + X'$.

The projection $\pi \circ \gamma$ is the standard rational normal curve γ' in \mathbb{C}^{m+p-1} . Similarly, the flag $K'_\bullet(s)$ osculating γ' at $\gamma'(s)$ is $\pi K_\bullet(s)$. Note that if L is a linear subspace of \mathbb{C}^{m+p} with $\gamma(\infty) \notin L$, then $\dim X \cap L = \dim \pi X \cap \pi L$. In particular, if $X \in \Omega_{J_m} K_\bullet(\infty)$, $s \neq \infty$, and $\alpha \in \binom{[m+p-1]}{p-1}$, then $\dim X' \cap K'_{(m+p-1)+1-\alpha_i} \geq (p-1) + 1 - i$ if and only if $\dim X \cap K_{m+p+1-(1+\alpha_i)} \geq p + 1 - (i+1)$. Thus we have

$$(19) \quad X \in \Omega_{J_m} K_\bullet(\infty) \cap \Omega_{\alpha^\perp} K_\bullet(s) \iff X' \in \Omega_\alpha K'_\bullet(s).$$

In fact, this induces an isomorphism of schemes.

This gives a strong equivalence between enumerative problems: If $\alpha^1, \dots, \alpha^n$ are in $\binom{[m+p-1]}{p-1}$ and s_1, \dots, s_n any complex numbers, then the map π induces an isomorphism between the schemes

$$\Omega_{J_m} K_\bullet(\infty) \cap \bigcap_{i=1}^n \Omega_{(\alpha^i)^\perp} K_\bullet(s_i) \quad \text{and} \quad \bigcap_{i=1}^n \Omega_{\alpha^i} K'_\bullet(s_i).$$

Part (ii) follows by noting that any real reparameterization of the rational normal curve γ induces an isomorphism of polynomial systems, thus preserves real solutions. Hence given $s_0, s_1, \dots, s_n \in \mathbb{P}_{\mathbb{R}}^1$, there is an equivalent system with $s_0 = \infty$.

It suffices to prove (iii) when β^\bullet covers α^\bullet in the partial order \preceq defined on Pieri Schubert data. Suppose Conjecture 3.1 fails for β^\bullet and β^\bullet covers α^\bullet with $\alpha^1 <_a \beta$ and $\alpha^2 = J_a$ as in Definition 3.6 (iii). Then there exist distinct real numbers s_1, s_3, \dots, s_n such that

$$(20) \quad \Omega_\beta K_\bullet(s_1) \cap \Omega_{\alpha^3} K_\bullet(s_3) \cap \dots \cap \Omega_{\alpha^n} K_\bullet(s_n)$$

is transverse with some complex p -planes in the intersection. We may assume without any loss that $s_1 = 0$. Then there is an open subset \mathcal{O} of the set of $(n-1)$ -tuples of real numbers s_3, \dots, s_n such that (20) is transverse and contains a complex p -plane X .

By the dimensional transversality results of [11], we may assume further that for $\beta' \in \binom{[m+p]}{p}$ and $(s_2, \dots, s_n) \in \mathcal{O}$, the intersection

$$\Omega_{\beta'} K_{\bullet}(0) \cap \bigcap_{i=3}^n \Omega_{\alpha^i} K_{\bullet}(s_i)$$

has the expected dimension and is transverse if 0-dimensional. This is empty if $|\beta'| > a + |\alpha^1|$, for dimension reasons. Thus

$$\left(\sum_{\alpha <_a \beta'} \Omega_{\beta'} K_{\bullet}(0) \right) \cap \bigcap_{i=3}^n \Omega_{\alpha^i} K_{\bullet}(s_i)$$

is transverse for $(s_3, \dots, s_n) \in \mathcal{O}$.

Fix $(s_3, \dots, s_n) \in \mathcal{O}$. By Proposition 3.5(i), there is an $\epsilon > 0$ such that for $|t| < \epsilon$

$$\Omega_{\alpha} K_{\bullet}(0) \cap \Omega_{J_a} K_{\bullet}(t) \cap \bigcap_{i=3}^n \Omega_{\alpha^i} K_{\bullet}(s_i)$$

is transverse. Here, when $t = 0$, replace $\Omega_{\alpha} K_{\bullet}(0) \cap \Omega_{J_a} K_{\bullet}(t)$ by $\sum_{\alpha <_a \beta} \Omega_{\beta'} K_{\bullet}(0)$. Since at $t = 0$ not all points in the intersection are real, the same holds for $0 < t < \epsilon$. But then Conjecture 3.1 fails for the Schubert data α^{\bullet} and completes the proof of Lemma 3.7. \blacksquare

3.4. An infinite family. We show that Conjecture 3.1 holds for an infinite family of nontrivial Schubert data.

Theorem 3.8. *Conjecture 3.1 holds for any m with $p = 2$ and Pieri Schubert data where one condition is J_{m-1} .*

Proof. By Lemma 3.7(iii), it suffices to show this for $\alpha^1 = \dots = \alpha^{m+1} = J_1$ and $\alpha^{m+1} = J_{m-1}$. Geometrically, we are looking for the 2-planes which meet a 2-plane and $m+1$ general m -planes nontrivially. We first show there are m such 2-planes. Let $L = K_2(\infty) = [0 \ I_2]$ and $M = K_m(0) = [I_m \ 0]$, and let $N_i = K_m(s_i)$, where s_1, \dots, s_m are distinct nonzero real numbers. For each one-dimensional subspace λ of L and each $1 \leq i \leq m$, the composition

$$M \hookrightarrow L \oplus M \simeq \mathbb{C}^{m+2} \twoheadrightarrow L \oplus M/(\lambda + N_i) \simeq \mathbb{C}$$

defines a linear form $\psi_{i,\lambda}$ on M . Each one-dimensional subspace μ of its kernel gives a 2-plane $\lambda \oplus \mu$ containing λ and meeting both M and N_i nontrivially.

Thus if X is a 2-plane meeting L, M , and each N_i nontrivially, then $H \cap L = \lambda$ and $H \cap M = \mu$ are lines with μ in the kernel of each form $\psi_{i,\lambda}$. Hence the forms are dependent. Similarly, if λ is a line in L such that the forms $\psi_{i,\lambda}$ are dependent, then any line μ they collectively annihilate gives a 2-plane $\lambda \oplus \mu$ meeting L, M , and each N_i nontrivially. It follows that the number of such 2-planes is the degree of the determinant of the forms $\psi_{i,\lambda}$, a polynomial in $\lambda \in \mathbb{P}(L) \simeq \mathbb{P}^1$. Since each form $\psi_{i,\lambda}$ is a linear function of λ , the determinant has degree m , so there are m 2-planes X meeting L, M , and each N_i nontrivially.

We compute this determinant and show it has only real roots. Let $\lambda = \lambda(x)$ be the span of the vector

$$(0, \dots, 0, 1, (m+1)x).$$

Let the rational normal curve γ have the parameterization

$$\gamma : s \longmapsto (1, -s, s^2, \dots, (-1)^{m+1} s^{m+1}).$$

Then $K_m(s)$, the osculating m -plane to γ at $\gamma(s)$, is the kernel of the matrix:

$$(21) \quad \begin{bmatrix} s^m & ms^{m-1} & \dots & \binom{m}{j}s^{m-j} & \dots & ms & 1 & 0 \\ 0 & s^m & \dots & \binom{m}{j-1}s^{m-j+1} & \dots & \binom{m}{2}s^2 & ms & 1 \end{bmatrix}.$$

If $R_j(s)$ is the linear form given by the j th row of this matrix, then

$$((m+1)x + ms_i)R_1(s_i) - R_2(s_i)$$

vanishes on $\lambda(x)$ and its restriction to M gives the form $\psi_{i,\lambda(x)}$. This restriction is represented by the vector $\Lambda(s_i, x)$ whose j th coordinate for $j = 0, \dots, m-1$ is:

$$\binom{m+1}{j} ((m-j+1)x s_i^{m-j} + (m-j)s_i^{m-j+1}).$$

We seek the determinant of the following matrix:

$$\begin{bmatrix} \Lambda(s_1, x) \\ \vdots \\ \Lambda(s_m, x) \end{bmatrix}.$$

This factors as $A \cdot B$, where A is the bidiagonal $m \times (m+1)$ -matrix

$$\begin{bmatrix} m & (m+1)x & 0 & & & \\ 0 & m^2 - 1 & m(m+1)x & & & \\ & & \ddots & \ddots & & \\ & 0 & \ddots & \binom{m+1}{j}(m-j) & \binom{m+1}{j}(m-j+1)x & 0 \\ & & & \ddots & \ddots & \\ & & 0 & & \binom{m+1}{2} & \binom{m+1}{2}2x \end{bmatrix}$$

and B is the $(m+1) \times m$ -matrix whose i, j th entry is s_j^{m+2-i} . Numbering the rows of A and the columns of B from 0 to m , we see that

$$\det(A(x) \cdot B) = \sum_{i=0}^m (-1)^i \det A_i(x) \det B_i,$$

where A_i is the matrix A with its i th column removed and B_i is the matrix B with its i th row removed. We find that

$$\begin{aligned} \det(A_i) &= m!(m+1-i)x^{m-i} \prod_{j=1}^m \binom{m+1}{j} \\ \det B_i &= e_i(s_1, \dots, s_m) s_1 s_2 \cdots s_m \cdot \prod_{j < k} (s_j - s_k) \end{aligned}$$

and so $\det(A \cdot B)$ is

$$m! \prod_{j < k} (s_j - s_k) \prod_{j=1}^m s_j^{\binom{m+1}{j}} \cdot \left(\sum_{i=0}^m (-1)^i (m-i+1) x^{m-i} e_i(s_1, \dots, s_m) \right).$$

Thus the coordinate x of the line λ satisfies the polynomial

$$P_m(s_1, \dots, s_m; x) := \sum_{i=0}^m (-1)^i (m-i+1) x^{m-i} e_i(s_1, \dots, s_m).$$

Since we have $e_i(s_1, \dots, s_m) = e_i(s_1, \dots, s_{m-1}) + s_m e_{i-1}(s_1, \dots, s_{m-1})$, we see that

$$P_m(s_1, \dots, s_m; x) = (x - s_m) P_{m-1}(s_1, \dots, s_{m-1}; x) + x \prod_{i=1}^{m-1} (x - s_i).$$

To complete the proof, we use induction to show that

$$(*) \quad \begin{aligned} &\text{If } 0 < s_1 < \dots < s_m, \text{ then the roots } r_1, \dots, r_m \text{ of } P_m \text{ satisfy:} \\ &0 < r_1 < s_1 < r_2 < s_2 < \dots < r_m < s_m. \end{aligned}$$

This suffices, if we can assume $0 < s_1 < \dots < s_m$. But we may assume this: Given a set of distinct real numbers $s_1, \dots, s_m, s_{m+1}, s_{m+2}$, we may assume $s_{m+2} = \infty$ and $s_{m+1} < s_1 < \dots < s_m$ and then apply the automorphism $s \mapsto s - s_{m+1}$ of $\mathbb{P}^1(\mathbb{R})$ which fixes $\infty = s_{m+2}$.

The case $m = 1$ of $(*)$ holds as $P_1(s_1; x) = 2x - s_1$. Suppose P_{m-1} satisfies $(*)$. Then the roots of $(x - s_m)P_{m-1}$ are $r_1 < r_2 < \dots < r_{m-1} < s_m$ and those of $x \prod_{i=1}^{m-1} (x - s_i)$ are $0 < s_1 < \dots < s_{m-1}$. Moreover the leading coefficients of both polynomials are positive. The result follows by the Intermediate Value Theorem: If $P(x)$ and $Q(x)$ are polynomials of degree n with positive leading coefficients and real interlaced roots p_i of P and q_i of Q

$$p_1 < q_1 < p_2 < q_2 < \dots < p_n < q_n,$$

then $P(x) + Q(x)$ has real roots r_i satisfying $p_i < r_i < q_i$, for $i = 1, \dots, n$. \blacksquare

3.5. Computational evidence. We have proven Conjecture 3.1 in a number of cases besides those of Theorem 3.8. We also have done many computations along the lines of those in Section 2.4. To describe these, we use the following compact notation. If a Schubert condition α is repeated k times in some Schubert data, we abbreviate that by α^k . Thus, the conditions of Conjecture 2.1 are written as J_1^{mp} .

Theorem 3.9. *Conjecture 3.1 holds for the following Schubert data.*

- (i) $(m, p) = (4, 2)$, $\alpha^* = J_2^4$. Here, $d(4, 2; J_2^4) = 3$.
- (ii) $(m, p) = (3, 3)$, $\alpha^* = J_2^4, J_1$. Here, $d(3, 3; J_2^4, J_1) = 3$.
- (iii) $(m, p) = (3, 3)$, $\alpha^* = (135)^2, J_1^3$. Here, $d(3, 3; (135)^2, J_1^3) = 6$.
- (iv) $(m, p) = (4, 3)$, $\alpha^* = 135^4$. Here, $d(4, 3; 135^4) = 8$.

Proof. We consider a polynomial system with parameters, give a universal eliminant, and show the eliminant has only real roots for distinct values of the parameters. We work in the local parameterization $\mathcal{X}_{\alpha^1, \alpha^2}$ of Section 3.2.

(i) Let $(m, p) = (4, 2)$ and $\alpha^\bullet = J_2^4$. The equations are

$$\text{maximal minors} \begin{bmatrix} 1 & s & s^2 & s^3 & s^4 & s^5 \\ 0 & 1 & 2s & 3s^2 & 4s^3 & 5s^4 \\ 0 & 0 & 1 & 3s & 6s^2 & 10s^3 \\ 1 & x_{12} & x_{13} & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & x_{25} & x_{26} \end{bmatrix} = 0$$

and the same equations with t replacing s . The ideal of these polynomials contains the following univariate polynomial g of degree 3 = $d(4, 2, J_2^4)$.

$$25x_{12}^3 - 25x_{12}^2(s+t) + x_{12}(19st + 6s^2 + 6t^2) - 3(s^2t + st^2)$$

whose discriminant has primitive part

$$9(s-t)^6 + 23s^2t^2(s-t)^2 + 9(s^6 + t^6).$$

Since $g(x_{12}; 1, 2)$ has roots

$$1, 1 \pm \frac{1}{5}\sqrt{7},$$

we have shown that g always has real roots, when s and t are distinct.

(ii) Let $m = p = 3$ and $\alpha^\bullet = J_2^4, J_1$. Here, \mathcal{X}_{J_2, J_1} consists of all matrices X of the form

$$\begin{bmatrix} 1 & x_{12} & 0 & 0 & 0 & 0 \\ 0 & 1 & x_{23} & x_{24} & x_{25} & 0 \\ 0 & 0 & 0 & 0 & 1 & x_{36} \end{bmatrix}$$

and our equations are

$$\det \begin{bmatrix} K_3(s) \\ X \end{bmatrix} = \text{maximal minors} \begin{bmatrix} K_2(t) \\ X \end{bmatrix} = 0$$

and the same equations with u replacing t . The ideal of these polynomials contains the following univariate polynomial g , here $e_1 = t + u$ and $e_2 = tu$.

$$\begin{aligned} x_{36}^3 - x_{36}(3s + 4e_1) + x_{36}(4e_1^2 + 3e_2 + 10se_1) - (6e_1e_2 + 8se_1^2 + se_2) \\ = (x_{36} - 2e_1)(x_{36}^2 - 2e_1x_{36} + 3e_2) - s(x_{36}^2 - 10e_1x_{36} + 8e_1^2 + e_2). \end{aligned}$$

These last two polynomials have roots

$$e_1 \pm \sqrt{e_1^2 - 3e_2}, \quad 2e_1 \quad \text{and} \quad \frac{5}{3}e_1 \pm \frac{\sqrt{e_1^2 - 3e_2}}{3},$$

which are interlaced. For example, if $e_1 > 0$, then

$$e_1 - \sqrt{e_1^2 - 3e_2} < \frac{5}{3}e_1 - \frac{\sqrt{e_1^2 - 3e_2}}{3} < e_1 + \sqrt{e_1^2 - 3e_2} < \frac{5}{3}e_1 + \frac{\sqrt{e_1^2 - 3e_2}}{3} < 2e_1.$$

When s, t, u are distinct and different from 0, g always has 3 real roots, by the Intermediate Value Theorem. We could also note that the discriminant of g

$$s^2(t-u)^4 + t^4(s-u)^2 + u^4(s-t)^2 + s^2t^2(s-t)^2 + s^2u^2(s-u)^2 + \\ (s-t)^2(s-u)^2(t-u)^2 + \frac{7}{2}(s^4(t-u)^2 + t^2(s-u)^4 + u^2(s-t)^4 + t^2u^2(t-u)^2)$$

is a sum of squares and $g(x_{36}; 1, 2, 3)$ has (floating point) roots

$$4.736, 7.756, 10.508.$$

(iii) Let $(m, p) = (3, 3)$ and $\alpha^\bullet = (135)^2, J_1^3$. Here, $\mathcal{X}_{135,135}$ consists of all matrices X of the form

$$\begin{bmatrix} 1 & x_{12} & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & x_{24} & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & x_{36} \end{bmatrix}$$

and our equations are

$$\det \begin{bmatrix} K_3(s) \\ X \end{bmatrix} = \det \begin{bmatrix} K_3(t) \\ X \end{bmatrix} = \det \begin{bmatrix} K_3(u) \\ X \end{bmatrix} = 0.$$

We write the universal eliminant, $g(x_{36})$, in terms of the elementary symmetric polynomials in s, t, u :

$$9x_{36}^6 - 48e_1x_{36}^5 + (64e_1^2 + 108e_2)x_{36}^4 - (288e_1e_2 - 198e_3)x_{36}^3 \\ + (320e_2^2 + 540e_1e_3)x_{36}^2 - 1200e_2e_3x_{36} + 1125e_3^2.$$

Evaluating the parameters (s, t, u) at $(1, 2, 3)$, we see that $g(x_{36}; 1, 2, 3)$ has roots:

$$1.491, 1.683, 3.210, 5.630, 9.213, 10.773.$$

The discriminant of g is a sum of squares. The primitive part of the discriminant is

$$e_3^4(4e_2^2e_1^2 - 15e_3e_1^3 - 15e_2^3 + 63e_3e_2e_1 - 81e_3^2)(256e_2^2e_1^2 - 768e_3e_1^3 - 768e_2^3 + 2592e_3e_2e_1 - 2187e_3^2)^2.$$

The second factor is a sum of squares

$$\frac{7}{2}(s-t)^2(s-u)^2(t-u)^2 + \frac{1}{2}s^2((t-u)^4 + t^2(s-u)^4 + u^2(s-t)^4)$$

Interestingly, the last (squared) factor is itself a sum of squares:

$$112(s-t)^2(u^4 + s^2t^2) + 112(t-u)^2(s^4 + t^2u^2) + 112(u-s)^2(t^4 + s^2u^2) + \\ 16(s-t)^2(s-u)^2(t-u)^2 + 309s^2t^2u^2 + 16(s^4(t^2 + u^2) + t^4(s^2 + u^2) + u^4(t^2 + u^2)).$$

(iv) Let $(m, p) = (4, 3)$ and $\alpha^\bullet = (135)^4$. Here, $\mathcal{X}_{135,135}$ consists of all matrices X of the form

$$\begin{bmatrix} 1 & x_{12} & x_{13} & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & x_{24} & x_{25} & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & x_{36} & x_{37} \end{bmatrix}$$

and our equations are

$$\text{maximal minors} \begin{bmatrix} K_3(s) \\ X \end{bmatrix} = \text{maximal minors} \begin{bmatrix} K_5(s) \\ X \end{bmatrix} = 0$$

and the same equations with t replacing s . In this case, the universal eliminant has 4 quadratic factors:

$$(36x_{12}^2 - x_{12}(12t + 30s) + 6st + 5s^2), \quad (36x_{12}^2 - x_{12}(12s + 30t) + 6st + 5t^2),$$

$$(3x_{12}^2 - 2x_{12}(s + t) + st), \quad \text{and} \quad (36x_{12}^2 - 30x_{12}(s + t) + 5t^2 + 14st + 5s^2).$$

When $s \neq t$ and neither is zero, we see that each has 2 real roots. \blacksquare

Observe that in all 4 cases, the discriminant was a sum of squares and the eliminant has the correct number of real roots for distinct values of the parameters. Of particular note is that the system in (ii) was not symmetric in the parameters and the Schubert data of (iv) was not Pieri Schubert data.

Table 3 gives the number of instances of Conjecture 3.1 we have checked.

α^\bullet :	$(J_2)^5$	$(J_2)^6$	$(J_2)^7$	$(J_2)^6$	$(J_3)^5$	$(135)^5$	$(135)^2, (J_1)^6$
m, p :	5,2	6,2	7,2	5,3	4,3	5,3	6,3
$d(m, p; \alpha^\bullet)$:	6	15	36	6	16	32	61
# checked	10,000	2821	504	10160	2002	400	294

TABLE 3. General Schubert data tested

4. TOTAL POSITIVITY

Previous sections have dealt with Schubert conditions given by flags osculating a real rational normal curve. Recently, Shapiro and Shapiro have conjectured that a generalization of this choice involving totally positive real matrices would also give only real solutions. We describe that here, prove the first nontrivial instance, and present some computational evidence in support of this generalization.

A real upper triangular matrix g with 1's on its diagonal is *totally positive* if every minor of g is positive, except those minors which vanish on all upper triangular matrices. Let \mathcal{TP} be the set of all totally positive, a multiplicative semigroup. Define a partial order on real flags F_\bullet by $F_\bullet < gF_\bullet$ if $g \in \mathcal{TP}$.

Conjecture 4.1. *For any $m, p > 1$, let α^\bullet be Schubert data for $\text{Grass}(p, m + p)$. If $F_\bullet^1 < \cdots < F_\bullet^n$ are real flags, then the Schubert varieties $\Omega_{\alpha^1} F_\bullet^1, \dots, \Omega_{\alpha^n} F_\bullet^n$ intersect transversally, with all points of intersection real.*

We will prove Conjecture 4.1 in the first nontrivial case of $m = p = 2$. First, we relate Conjecture 4.1 to Conjecture 3.1. Let $K_\bullet(s)$ be the square matrix of size $(m + p)$ whose i, j th entry is $\binom{j-i}{i-1} s^{j-i}$ (cf. (1)). If $s > 0$, then $K_\bullet(s)$ is totally positive and for any s, t we have $K_\bullet(s) \cdot K_\bullet(t) = K_\bullet(s + t)$. To see this, first recall that \mathcal{TP} is generated as a semigroup by $\exp(E_{i,i+1})$, where $E_{i,i+1}$ is the elementary matrix whose only non-zero entry is in position $i, i + 1$ [27]. These assertions follow from the observation that

$$K_\bullet(s) = \exp(sN),$$

where N is the nilpotent matrix whose only non-zero entries are $(1, 2, \dots, m+p-1)$ lying just above its main diagonal.

Theorem 3.3 holds in this new setting. For this, we alter the notion of Pieri Schubert data α^\bullet to Schubert data $\alpha^1, \dots, \alpha^n$ where all except possibly α^1 and α^n are Pieri conditions.

Theorem 4.2. *Let $a, b > 1$ and suppose that Conjecture 4.1 holds for $(m, p) = (a, b)$ and Schubert data $\alpha^\bullet = (J_1)^{mp}$. Then Conjecture 4.1 holds for any Pieri Schubert data for $\text{Grass}(p, m+p)$ where $(m, p) \leq (a, b)$ or (b, a) coordinatewise.*

Proof. The arguments used to prove Theorem 3.3 work here with minor adjustments.

We first remark that total positivity, and hence our order $<$ on real flags, is defined with respect to a choice of ordered basis for \mathbb{R}^{m+p} . Suppose that e_1, \dots, e_{m+p} is the basis we used to define this order. Then $F_\bullet < G_\bullet$ is and only if $G_\bullet <' F_\bullet$, where $<'$ is defined with respect to the basis $e_1, -e_2, e_3, -e_4, \dots$. Similarly, if the basis e_1, \dots, e_{m+p} is orthonormal, then $F_\bullet < G_\bullet$ if and only if $F_\bullet^\perp <' G_\bullet^\perp$, where $<'$ is defined with respect to the basis in reverse order e_{m+p}, \dots, e_2, e_1 . Thus

$$F_\bullet^1 < F_\bullet^2 < \dots < F_\bullet^n \iff F_\bullet^n <' \dots <' F_\bullet^2 <' F_\bullet^1$$

so that Conjecture 4.1 holds for Schubert data α^\bullet if and only if it holds for the data in reverse order. (This is the only rearrangement we used in the proof of Lemma 3.7.) Similarly, the analogue of Lemma 3.7(i) holds. For the analogue of Lemma 3.7(ii), permute the last two Schubert conditions, so that β^\bullet is still Pieri Schubert data, in our new, restricted definition.

Finally, in the proof of Lemma 3.7(iii), replace s_3, \dots, s_n in defining the set \mathcal{O} by fixing F_\bullet^1 to be the standard flag represented by the matrix I_{m+p} and let \mathcal{O} be the set of all

$$F_\bullet^1 < \dots < F_\bullet^n$$

where the appropriate transversality conditions hold. Since \mathcal{TP} is open, it follows that there exists $\epsilon > 0$ and totally positive matrix M (which stabilizes F_\bullet^1) such that if $0 < s < \epsilon$, then $F_\bullet^1 < M \cdot K_\bullet(s) \cdot F_\bullet^1 < F_\bullet^2$. Then the same arguments used to prove Theorem 3.3 suffice. In particular, the analog of Proposition 3.2 also holds in this setting. \blacksquare

Totally positive matrices have a useful description. Let \mathcal{U} be the group of real unipotent (upper triangular) matrices. Then \mathcal{TP} is a connected component of the complement of a hypersurface H in \mathcal{U} defined by the vanishing of all minors consisting of the first i rows and last i columns [38]. This has a geometric description.

Associating a matrix to a flag as in Section 3.2, we may identify \mathcal{U} with a Zariski open subset of the real flag manifold. Then the hypersurface H is the union of all positive codimension Schubert varieties defined by the flag determined by the identity matrix.

Given a matrix $M \in \mathcal{U}$, the translate $\mathcal{TP} \cdot M$ is a component of the complement of all Schubert varieties of positive codimension defined by the flag given by M . Similarly, given a totally positive matrix M , the set of upper triangular matrices N

for which there exists a totally positive g with $gN = M$ is the component of this complement containing the identity matrix.

Let $F_{\bullet}^1 < \dots < F_{\bullet}^n$ be real flags. Using a real automorphism of the flag manifold, we may assume that $F_{\bullet}^1 = K_{\bullet}(0) = I_{m+p}$. Then $F_{\bullet}^2, \dots, F_{\bullet}^n \in \mathcal{TP}$, since they are all translates of the identity by totally positive matrices. Also, $F_{\bullet}^1, \dots, F_{\bullet}^{n-1}$ are in the same component of the complement of all positive dimensional Schubert cells defined by F_{\bullet}^n . If we now consider a real coordinate transformation fixing F_{\bullet}^1 , but with F_{\bullet}^n becoming $K_{\bullet}(\infty)$, then this complement becomes \mathcal{TP} , in these new coordinates.

Thus we may work in the local coordinates $\mathcal{X} := \mathcal{X}_{\alpha^1, \alpha^n}$. We do this in our proof of the following theorem and in subsequent calculations.

Theorem 4.3. *Conjecture 4.1 holds for $m = p = 2$ and Schubert data $(J_1)^4$.*

Proof. Let $F, G \in \mathcal{TP}$ be totally positive matrices and set $H = G \cdot F$. When $m = p = 2$, $\mathcal{X} = \mathcal{X}_{J_1, J_1}$ is the set of matrices

$$\begin{bmatrix} 1 & a & 0 & 0 \\ 0 & 0 & 1 & b \end{bmatrix}.$$

For a matrix L , let L_{ij} denote the 2×2 -minor of L given by the first two rows and columns i and j . Then the equations for a 2-plane in \mathcal{X} to meet the flags given by F and H are

$$\begin{aligned} f &:= F_{24} - bF_{23} - aF_{14} + abF_{13} \\ h &:= H_{24} - bH_{23} - aH_{14} + abH_{13} \end{aligned}$$

The lexicographic Gröbner basis for this (with $a < b$) is

$$\begin{aligned} H_{13}f - F_{13}h &= J_{14} - bJ_{24} - aJ_{34} \\ (H_{14} - bH_{13})f - (F_{14} - bF_{13})h &= J_{13} - b(J_{23} + J_{14}) + b^2J_{24}, \end{aligned}$$

where J_{ij} is the ij th minor of the matrix:

$$\begin{bmatrix} F_{24} & F_{23} & F_{14} & F_{13} \\ H_{24} & H_{23} & H_{14} & H_{13} \end{bmatrix}.$$

We may write the the discriminant of the quadratic equation for b as follows

$$(J_{23} + J_{14})^2 - 4J_{13}J_{24} = (L_{23} + L_{14})^2 - 4L_{13}L_{24},$$

where L is the matrix

$$\begin{bmatrix} F_{13} & F_{14} & H_{13} & H_{14} \\ F_{23} & F_{24} & H_{23} & H_{24} \end{bmatrix}.$$

Thus we will have two real roots for our original system if and only if

$$\Lambda(B) := L_{13} - B(L_{23} + L_{14}) + B^2L_{24} = 0$$

has 2 real solutions. Painstaking calculations reveal that $\Lambda(1) = -G_{12}G_{34} < 0$. Since $L_{24} = H_{13}H_{24} - H_{23}H_{14} = H_{12}H_{34}$ by the Plücker relations, we see that $L_{24} > 0$ and so $\Lambda(B) = 0$ will have 2 real solutions. \blacksquare

Table 4 shows the number of instances of Conjecture 4.1 that we have verified.

α^\bullet	$(J_1)^6$	$(J_2)^5$	$(135)^4$	$(J_1)^8$	$(J_2)^6$	$(135)(136)(J_1)^5$
m, p	3,2	5,2	4,3	4,2	6,2	4,3
d	5	6	8	14	15	25
# checked	12000	4000	4000	1500	300	150

TABLE 4. Instances checked

5. FURTHER REMARKS

We present a counterexample to the original conjecture of Shapiro and Shapiro and close with a discussion of further questions.

5.1. A counterexample to the original conjecture. The original conjecture of Shapiro and Shapiro concerned the M -property for flag manifolds [37]. An algebraic set X defined over \mathbb{R} has the M -property if the sum of the $\mathbb{Z}/2\mathbb{Z}$ -Betti numbers of $X(\mathbb{R})$ and of $X(\mathbb{C})$ are equal. Shapiro and Shapiro conjectured that an intersection of Schubert cells in a flag manifold has the M -property, if the cells are defined by flags osculating the rational normal curve at real points. When such an intersection is zero-dimensional all of its points are real. It is this consequence we have been studying.

While there is much evidence in support of this conjecture for zero dimensional intersections in a Grassmannian (Conjectures 2.1, 3.1, and 4.1), it does not hold for more general flag manifolds. In fact, we give a counter example in the simplest enumerative problem in a flag manifold that does not reduce to an enumerative problem in a Grassmannian.

Counterexample 5.1. Consider the manifold $\mathbb{F}(2, 3; 5)$ consisting of partial flags $X \subset Y$ in \mathbb{C}^5 with $\dim X = 2$ and $\dim Y = 3$. This manifold has dimension 8; the projection to $Grass(2, 5)$ has fibre over a 2-plane X equal to $\mathbb{P}(\mathbb{C}^5/X) \simeq \mathbb{P}^2$. Given general 2-planes a, b , and c and general 3-planes A, B , and C , there are 4 flags $X \subset Y$ which satisfy

$$(22) \quad \begin{array}{l} X \text{ meets } a, B, \text{ and } C \text{ nontrivially} \\ \dim Y \cap A \geq 2 \text{ and } Y \text{ meets } b \text{ and } c \text{ nontrivially} \end{array}$$

That this number is 4 may be verified using the Schubert calculus for a flag manifold [17] or the equations we give below.

Let $K_\bullet(s)$ be the flag of subspaces osculating the standard rational normal curve. Set

$$\begin{array}{ll} a := K_2(4) & A := K_3(0) \\ b := K_2(1) & B := K_3(3) \\ c := K_2(-5) & C := K_3(-1) \end{array}$$

We claim that of the 4 flags $X \subset Y$ satisfying (22) for this choice of a, b, c, A, B, C , 2 are real and 2 are complex.

We outline the computation. Choose local coordinates for $\mathbb{F}(2, 3; 5)$ as follows. Let Y be the row space of the 3×5 -matrix

$$\begin{bmatrix} 0 & 0 & 1 & x_{14} & x_{15} \\ 1 & 0 & x_{23} & x_{24} & x_{25} \\ 0 & 1 & x_{33} & x_{34} & x_{35} \end{bmatrix}$$

and X be the row space of its last 2 rows. We seek the solutions to the following overdetermined system of polynomials:

$$\det \begin{bmatrix} K_2(1) \\ Y \end{bmatrix} = \det \begin{bmatrix} K_2(-5) \\ Y \end{bmatrix} = \det \begin{bmatrix} K_3(3) \\ X \end{bmatrix} = \det \begin{bmatrix} K_3(-1) \\ X \end{bmatrix} =$$

$$\text{maximal minors} \begin{bmatrix} K_2(4) \\ X \end{bmatrix} = \text{maximal minors} \begin{bmatrix} K_3(0) \\ Y \end{bmatrix} = 0$$

These polynomials generate a zero-dimensional ideal containing the following univariate polynomial, which is part of a lexicographic Gröbner basis satisfying the Shape Lemma:

$$27063 - 117556x_{14} - 5952x_{14}^2 - 10416x_{14}^3 + 32400x_{14}^4.$$

This has roots

$$-.736 \pm 1.30\sqrt{-1}, \quad .227, \quad 1.62.$$

Thus 2 of the flags are complex.

5.2. Further questions. While Counterexample 5.1 shows that we cannot guarantee all points of intersection real when the Schubert varieties are given by flags osculating a real rational normal curve, a number of questions remain (besides the resolution of the conjectures of the previous sections). There remains the original question of Fulton:

Question 1: Given Schubert data for a flag manifold, do there exist real flags in general position whose corresponding Schubert varieties have *only* real points of intersection?

In every case we know, this does happen. For instance, if we change the 3-plane B to $K_3(2)$ in Counterexample 5.1, then all 4 solution flags are real. There is also the following result, showing this holds in infinitely many cases. Let $\mathbb{F}(2, n-2; n)$ be the manifold of flags $X \subset Y$ in \mathbb{C}^n where $\dim X = 2$ and $\dim Y = n-2$. A Grassmannian Schubert condition is a Schubert condition on the flag $X \subset Y$ which only imposes conditions on one of the subspaces. We likewise define Grassmannian Schubert data. For example, Counterexample 5.1 involves Grassmannian Schubert data.

Proposition 5.2 (Theorem 13 of [41]). *Given any Grassmannian Schubert data for $\mathbb{F}(2, n-2; n)$, there exist real flags whose corresponding Schubert varieties meet transversally with all points of intersection real.*

The beauty of the conjectures of Shapiro and Shapiro is that they give a simple algorithm for selecting the flags defining the Schubert varieties.

Question 2: Can the choice of flags in Question 1 (or Proposition 5.2) be made effective? In particular, is there an algorithm for selecting these flags?

While computing the examples described here, we have made a number of observations which deserve further scrutiny. These concern eliminant polynomials in the ideals defining the intersections of Schubert varieties in the local coordinates we have been using.

Suppose we have Schubert data α^\bullet , and have chosen local coordinates either for the Grassmannian or are working in $\mathcal{X}_{\alpha^n, \alpha^{n-1}}$. Conjecture 3.1 or 4.1 may be formulated in terms of a parameterized system of polynomials with parameters either s_1, \dots, s_n in the case of Conjecture 3.1 or $(n-1)$ -tuples of totally positive matrices (or in terms of some parameterization of \mathcal{TP} [2]). For each of the coordinates, the ideal of this system contains a universal eliminant, which is the minimal univariate polynomial in that coordinate with coefficients rational functions in the parameters.

We ask the following questions about the eliminant.

Question 3: Does the universal eliminant have degree equal to the generic number of solutions? That is, do generic solutions satisfy the shape lemma?

Question 4: Let Δ be the discriminant of the polynomial system, a polynomial in the parameters which vanishes when there are solutions with multiplicities.

- a) Is the locus $\Delta \neq 0$ connected?
- b) In the case of Conjecture 3.1, where Δ is a polynomial in the parameters s_1, \dots, s_n , is Δ always a sum of squares of polynomials?
- c) If so, are these polynomials monomials in the s_i and their differences $(s_i - s_j)$? This would imply that the polynomial systems are always multiplicity-free for distinct real values of the parameters, and hence the stronger version of Theorem 3.3 mentioned in Remark 3.4.

The discriminants we have computed for instances of the conjectures for the Grassmannian (including the discriminant for system of Theorem 4.3) are always non-negative when the parameters are distinct. For the case of Counterexample 5.1, we computed a discriminant for a simpler, but equivalent system, in the spirit of sections 2.5 and 3.2. This polynomial in parameters s_1, s_2, t_1, t_2 is symmetric in the s 's and in the t 's separately (and in the transformation $s_i \leftrightarrow t_i$) and has degree 24. It has three factors, the first of degree 20 with 857 terms, and the square

$$(2s_1s_2 + 2t_1t_2 - (s_1 + s_2)(t_1 + t_2))^2.$$

While this factor will not prevent the discriminant from being a sum of squares, this factor shows that there is a choice of distinct parameters for which the discriminant vanishes. Indeed, if we set $s_1 = 3, s_2 = 6, t_1 = 9$, and $t_2 = 5$, then this factor vanishes, and the resulting system has a root of multiplicity 2. This also explains why different values of the parameters in Counterexample 5.1 give different numbers of real and complex solutions.

Question 5: When the universal eliminant factors over \mathbb{Z} , it reflects either some underlying geometry or some interesting arithmetic. More generally, one might ask about the Galois group of these enumerative problems [20], or the Galois group of

the universal eliminant. For instance, is it the full symmetric group? That is not always the case, as the example of Theorem 3.9(iv) shows.

Question 6: In many cases with the substitution of $s_i = i$, the eliminant factors over the integers. This happens in Conjecture 1.1, Theorem 2.3, Theorem 3.9(i) and (iv), and in other cases. Table 5 lists the degrees of the factors in the case of Conjecture 1.1. Why does this choice of $s_i = i$ induce a factorization? Is there any

m, p	3,2	4,2	5,2	6,2	7,2	3,3	3,4
$d_{m,p}$	5	14	42	132	429	42	462
Factors	2,3	6,8	10,32	20,112		6,36	16,30,416

TABLE 5. Factorization of the eliminant

special geometry or interesting arithmetic here? If 2 parameters are allowed to come together, then the resulting ideal factors in a way respecting the product of Schubert classes, by the Corollary to Theorem 1 in [12]. From the Schubert calculus, we would expect factors of 9 and 5 for $(m, p) = (2, 4)$, 14 and 28 for $(m, p) = (2, 5)$, and 21 and 21 for $(m, p) = (3, 3)$, but these do not appear in Table 5.

ACKNOWLEDGEMENTS

I thank Boris Shapiro for sharing these conjectures and Joachim Rosenthal with whom I found the first strong evidence for their validity [33]. Bernd Sturmfels helped me learn computer algebra and suggested using discriminants. This work also benefitted from discussions with David Eisenbud, Ioannis Emiris, Jean-Charles Faugère, Bill Fulton, Birk Huber, Bernard Mourrain, Bruce Reznick, Fabrice Rouillier, Jan Verschelde, and Paul Zimmermann. Lastly, this would not be possible were it not for the authors of the following software packages: Linux, Maple, Singular, Macaulay2, MuPAD, FGB, RealSolving, PHC, and also my brothers Larry Sottile and Joe Sottile who helped me build the computer which performed most of these experiments.

REFERENCES

- [1] T. ANDO, *Totally positive matrices*, Lin. Alg. Appl., 90 (1987), pp. 165–219.
- [2] A. BERENSTEIN, S. FOMIN, AND A. ZELEVINSKY, *Parameterizations of canonical bases and totally positive matrices*. 1995.
- [3] R. W. BROCKETT AND C. I. BYRNES, *Multivariable Nyquist criteria, root loci and pole placement: A geometric viewpoint*, IEEE Trans. Automat. Control., AC-26 (1981), pp. 271–284.
- [4] C. I. BYRNES, *Algebraic and geometric aspects of the control of linear systems*, in Geometric Methods in Linear systems Theory, C. I. Byrnes and C. F. Martin, eds., D. Reidel, Dordrecht, 1980, pp. 85–124.
- [5] C. I. BYRNES, *Control theory, inverse spectral problems, and real algebraic geometry*, in Differential Geometric Methods in Control Theory, R. W. Brockett, R. S. Millman, and H. J. Sussmann, eds., Birkhäuser, Boston, 1982, pp. 192–208.
- [6] C. I. BYRNES, *Pole assignment by output feedback*, in Three Decades of Mathematical Systems Theory, H. Nijmeijer and J. M. Schumacher, eds., vol. 135 of Lecture Notes in Control and Inform. Sci., Springer-Verlag, Berlin, 1989, pp. 31–78.

- [7] C. I. BYRNES AND P. K. STEVENS, *Global properties of the root-locus map*, in Feedback Control of Linear and Non-Linear Systems, D. Hinrichsen and A. Isidori, eds., vol. 39 of Lecture Notes in Control and Inform. Sci., Springer-Verlag, Berlin, 1982.
- [8] M. CHOI, T. LAM, AND B. REZNICK, *Even symmetric sextics*, Math. Z., 195 (1987), pp. 559–580.
- [9] ———, *Sums of squares of real polynomials*, in *K-theory and algebraic geometry: connections with quadratic forms and division algebras* (Santa Barbara, 1992), vol. 58.2 of PSPUM, AMS, 1995, pp. 103–126.
- [10] D. COX, J. LITTLE, AND D. O'SHEA, *Using Algebraic geometry*, GTM 185, Springer-Verlag, New York, 1998.
- [11] D. EISENBUD AND J. HARRIS, *Divisors on general curves and cuspidal rational curves*, Invent. Math., 74 (1983), pp. 371–418.
- [12] ———, *When ramification points meet*, Invent. Math., 87 (1987), pp. 485–493.
- [13] J.-C. FAUGÈRE, *FGB*. See <http://posso.lip6.fr/~jcf>.
- [14] J.-C. FAUGÈRE, P. GIANNI, D. LAZARD, AND T. MORA, *Efficient computation of zero-dimensional Groebner bases by change of ordering*, J. Symb. Comp., 16 (1993), pp. 329–344.
- [15] J.-C. FAUGÈRE, F. ROUILLIER, AND P. ZIMMERMANN, *Verifying the conjecture of Shapiro and Shapiro for $m, p = 3, 4$* , 1998.
- [16] W. FULTON, *Introduction to Intersection Theory in Algebraic Geometry*, CBMS 54, AMS, 1996. second edition.
- [17] ———, *Young Tableaux*, Cambridge University Press, 1996.
- [18] L. GONZALEZ-VEGA, F. ROUILLIER, M. ROY, AND G. TRUJILLO, *Symbolic recipes for real solutions*. to appear.
- [19] G.-M. GREUEL, G. PFISTER, AND H. SCHÖNEMANN, *Singular version 1.2 user manual*, Tech. Rep. 21, Centre for Computer Algebra, June 1998. <http://www.mathematik.uni-kl.de/~zca/Singular>
- [20] J. HARRIS, *Galois groups of enumerative problems*, Duke Mathematical Journal, 46 (1979), pp. 685–724.
- [21] D. HILBERT, *Über die Darstellung definiter Formen als Summe von Formen-quadrate*, Math. Ann., 32 (1888), pp. 342–350.
- [22] W. V. D. HODGE AND D. PEDOE, *Methods of Algebraic Geometry*, vol. II, Cambridge University Press, 1952.
- [23] B. HUBER, F. SOTTILE, AND B. STURMFELS, *Numerical Schubert calculus*. J. Symb. Comp., 26 (1998) 767–788.
- [24] S. KLEIMAN, *The transversality of a general translate*, Comp. Math., 28 (1974), pp. 287–297.
- [25] S. KLEIMAN AND D. LAKSOV, *Schubert calculus*, Amer. Math. Monthly, 79 (1972), pp. 1061–1082.
- [26] A. G. KOUCHNIRENKO, *A Newton polyhedron and the number of solutions of a system of k equations in k unknowns*, Usp. Math. Nauk., 30 (1975), pp. 266–267.
- [27] C. LOEWNER, *On totally positive matrices*, math. Zeitschr., 63 (1955), pp. 338–340.
- [28] C. F. MARTIN AND R. HERMANN, *Applications of algebraic geometry to system theory: The McMillan degree and Kronecker indices as topological and holomorphic invariants*, SIAM J. Control, 16 (1978), pp. 743–755.
- [29] A. S. MORSE, W. A. WOLOVICH, AND B. D. O. ANDERSON, *Generic pole assignment: Preliminary results*. Proc. 20th IEEE Conference on Decision and Control, San Diego, 1981.
- [30] M. S. RAVI, J. ROSENTHAL, AND U. HELMKE, *On output feedback invariants and cascade equivalence of systems*, in Proc. of the 36th IEEE Conference on Decision and Control, IEEE, 1997, pp. 4243–4248. San Diego, California.
- [31] F. RONGA, A. TOGNOLI, AND T. VUST, *The number of conics tangent to 5 given conics: the real case*, Rev. Mat. Univ. Complut. Madrid, 10 (1997), pp. 391–421.
- [32] J. ROSENTHAL, J. M. SCHUMACHER, AND J. C. WILLEMS, *Generic eigenvalue assignment by memoryless real output feedback*, Systems & Control Lett., 26 (1995), pp. 253–260.

- [33] J. ROSENTHAL AND F. SOTTILE, *Some remarks on real and complex output feedback*, Systems & Control Lett., 33 (1998), pp. 73–80. For a description of the computational aspects, see <http://www.nd.edu/~rosen/pole/>.
- [34] F. ROUILLIER, *RealSolving*. Available at <http://www.loria.fr/~rouillie>
- [35] ———, *Solving zero-dimensional systems through the Rational Univariate Representation*. Projet Polka, INRIA Lorraine, 1998.
- [36] H. SCHUBERT, *Beziehungen zwischen den linearen Räumen auferlegbaren charakteristischen Bedingungen*, Math. Ann., 38 (1891), pp. 588–602.
- [37] B. SHAPIRO AND M. SHAPIRO, *The M-property of flag manifolds*, Top. and Appl., 43 (1992), pp. 65–81.
- [38] ———, *On the boundary of totally positive upper triangular matrices*, Lin. Alg. Appl., 231 (1995), pp. 105–109.
- [39] F. SOTTILE, *Enumerative geometry for real varieties*, in Algebraic Geometry, Santa Cruz 1995, J. Kollár, R. Lazarsfeld, and D. Morrison, eds., vol. 62, Part 1 of Proc. Sympos. Pure Math., Amer. Math. Soc., 1997, pp. 435–447.
- [40] ———, *Enumerative geometry for the real Grassmannian of lines in projective space*, Duke Math. J., 87 (1997), pp. 59–85.
- [41] ———, *Real enumerative geometry and effective algebraic equivalence*, J. Pure Appl. Alg., 117 & 118 (1997), pp. 601–615. Proc., MEGA'96.
- [42] ———, *The conjecture of Shapiro and Shapiro: summary and an archive of computations*. <http://www.math.wisc.edu/~sottile/pages/shapiro/index.html>, 1998.
- [43] S. A. STRØMME, *On parameterized rational curves in Grassmann varieties*, in Space Curves, F. Ghione, C. Peskine, and E. Sernesi, eds., vol. 1266 of Lecture Notes in Mathematics, Springer-Verlag, 1987, pp. 251–272.
- [44] B. STURMFELS, *Polynomial equations and convex polytopes*. Amer. Math. Monthly, to appear, 1998.
- [45] V. L. SYRMOS, C. T. ABDALLAH, P. DORATO, AND K. GRIGORIADIS, *Static output feedback—a survey*, Automatica J. IFAC, 33 (1997), pp. 125–137.
- [46] J. VERSCHelde, *PHC: A general-purpose solver for polynomial systems by homotopy continuation*, Tech. Rep. TW 265, Dept. of Computer Science, K.U.Leuven, 1997. <http://www.msri.org/people/members/jan/PHCpack/paper.html>
- [47] J. VERSCHelde, *Numerical evidence of a conjecture in real algebraic geometry*. MSRI preprint # 1998-064 1998.
- [48] J. C. WILLEMS AND W. H. HESSELINK, *Generic properties of the pole placement problem*, in Proc. of the 7th IFAC Congress, 1978, pp. 1725–1729.

FRANK SOTTILE, DEPARTMENT OF MATHEMATICS, UNIVERSITY OF WISCONSIN, VAN VLECK HALL, 480 LINCOLN DRIVE, MADISON, WISCONSIN 53706-1388, USA

E-mail address: sottile@math.wisc.edu

URL: <http://www.math.wisc.edu/~sottile>